Massive Collaborative Wireless Sensor Network Structure Based on Cloud Computing

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Abstract—To explore the wireless sensor network (WSN) structure, the cooperative WSN architecture of mass data processing based on cloud computing is studied. The technology of WSN and cloud computing is deeply discussed. The system and node structure of WSN are studied by theoretical analysis method, and the performance of the WSN is studied by using the numerical simulation method. The mass data processing technology based on Map Reduce and its application in WSN are discussed. The numerical simulation method is used to experiment on the architecture of SVC4WSN and MD4LWSN. The relationship between the optimal network number and the node communication radius at different node density is verified. Moreover, the energy and time delay Reduce path is compared with three protocols of LEACH, PEGASIS and PEDAP. The results show that the two Reduce paths have better performance in both network survival time and the total time slot of data acquisition.

Keywords—WSN; massive data processing; cloud computing

1 Introduction

Since twenty-first Century, with the rapid popularization of automatic information generation equipment represented by sensors and intelligent recognition terminals, people can accurately perceive the data of the physical world in real time. As a new component of information technology, Internet of Things (IoT) has been widely concerned by many countries all over the world. The United States first regards the IoT as the direction of national strategic development. In 2009, International Business Machine Company put forward the strategy of "intelligent earth". In the same year, the Chinese government established the central goal of building a "perceived China" based on the IoT. According to the Forrester of the authoritative advisory body, by 2020, the world's interconnected business will reach 30:1 compared to the human communication business; by 2035, the wireless sensor network (WSN) terminal of China will reach hundreds of billions; and by 2050, the sensor will be ubiquitous in life and this is the scale effect of intelligent devices in IoT.

The IoT depicts a new way of interaction between human beings and the physical world, and its grand blueprint has been widely accepted. IoT, that is, the Internet connected by objects, mainly contains two meanings: the first is that its core and foundation
is still the Internet, and it is the network extended on the basis of the Internet; and the second is that its user end is extended to information exchange and communication between any object. At present, the mainstream definition of the IoT is proposed by the International Telecommunication Union in 2005. That is, the IoT is a network that connects any objects with the Internet for information exchange and communication to realize intelligent identification, location, tracking, monitoring and management of objects through information sensing devices, such as radio frequency identification, infrared sensors, global positioning systems, and laser scanners. Compared to the traditional communication network or the Internet, the IoT will make every object in the world corresponding to a unique code, collect all kinds of signals through information communication and network technology, convert it into information flow, and combine it with the Internet, so as to form a new communication between people and objects and between objects and objects.

The IoT has become the next high point in the field of information technology. In the face of the coming large-scale application, the data processing and architecture of the IoT have become the focus of researches. The collaborative WSN architecture based on distributed mass data processing is proposed in the face of this urgent need. With the large-scale development of WSNs used by the perceptive layer in the IoT, the massive data produced in this network need to be processed in time, and users expect to get more useful information from these data. Therefore, the research on massive data processing and collaborative WSN architecture has realistic and long-term significance.

2 Literature review

IoT is recognized as one of the key technologies for the third breakthroughs in the IT industry after the Internet and mobile communications technology. The first and foremost problem in the construction and application of IoT is how to deal with complex data. All the countries in the world have paid great attention to WSN and its applications. The informational strategies such as the “intelligent earth” of the United States, the “u-Japan” of Japan, the “u-Korea” of Korea, the “European digital plan” of the EU and China’s “perceived China” have developed rapidly. The common points of these informational strategies include the integration of various information technologies, the breakthrough of Internet limitations, the access of objects to information network and the realization of object networking. On the basis of the Internet popularization, the information technology is applied to various fields, which has far-reaching influence on all aspects of the national economy and social life. The development of the future information industry extends and breaks through from the information network to the overall perception and intelligent application.

The United States is the birthplace of WSNs, and related technologies are in the leading position in the world. City SENSE, developed by the Harvard University and BBN Company, is used to monitor urban weather and environmental pollution. The goal is to build the world’s first wireless network that can deliver sensor data in real time throughout the city. The earth central system is a project being developed by the Hewlett-Packard Company and will be completed this year. It aims at creating the
mathematical and physical basis of the technology needed to create a new information ecosystem, including one trillion nanoscale sensors and drives embedded in the large environment. The computer systems, software and services are connected through a variety of networks to exchange information between the analysis engine, the storage system and the end users.

The European Union (EU) has also carried out a lot of researches and applications in WSNs. On September 15, 2009, the EU issued the "road map of the strategic research on the IoT", and put forward the research field and research route of the EU's IoT for 2010-2020 years. This road map indicates that all "objects" in the vision of the IoT have a common feature, that is, "they can add sensors to interact with the environment where they are". The future technological development assumptions and research needs are put forward, among which sensors and WSNs are involved. The main research projects of WSNs in the EU include ARTEMIS (Advanced Research and Technology for Embedded Intelligence and Systems). ARTEMIS studied new multidisciplinary coordination and control principles related to large-scale WSNs and actuator networks, including integrated control, computing and communications (C3) strategy.

In the distributed large-scale WSN data processing, many scholars at home and abroad have put forward corresponding solutions.

Arora et al. (2016) proposed a data storage method based on distributed space and time similarity to improve the efficiency of data query in WSNs [1]. Zahurul et al. (2016) allocated the tasks of sensor nodes and put forward a kind of energy efficient data acquisition scheme with the aid of the correlation between space and time [2]. Olofsson et al. (2016) put forward the protocol for large scale WSN data collection to achieve the goal of maximizing the network lifetime by balancing the energy consumption of the whole network [3]. Wei et al. (2017) proposed a new storage and range query scheme in the data of the sensor network, reducing the message cost and obtaining the load balancing in the network [4]. Khasawneh et al. (2017) proposed the distributed data collection and distributed data aggregation algorithm in the large-scale heterogeneous WSN under the generalized physical interference model [5]. Kumar and Kumar (2016) proposed a method based on distributed reinforcement learning routing in large-scale WSNs, and obtained satisfactory end-to-end delay and reliability [6]. Mai et al. (2016) put forward a distributed algorithm based on time and energy tradeoff for aggregation of large sensor network data [7].

Gasparri et al. (2016) put forward the combination of WSN and cloud computing technology for telemedicine data acquisition and monitoring. The data transferred from sensors are stored and processed in the cloud. The platform-as-a-service (PaaS) and software-as-a-service technologies are adopted in the cloud-based architecture. The proposed architecture can provide real-time data acquisition and reduce acquisition errors, which simplifies the deployment process [8]. Ghari et al. (2016) proposed the application and development technology of WSN based on the PaaS platform. The sensor nodes can be easily developed by using this technology. With the help of the tools provided by the PaaS platform, users can quickly develop the business application which combines sensor data [9]. Zhang et al. (2017) put forward a sensor network environment monitoring application architecture model based on cloud computing [10].
the help of the elastic computing characteristics of Amazon EC2, the dynamic load requirements of the actual environmental application are met.

In general, the existing schemes cannot effectively solve the heterogeneity of the sensor network data, the support for the current mass data processing of the sensor network is not enough, and the data collected cannot be used very well. In view of this, the new WSN is combined with cloud computing architecture technology. It is mainly divided into two aspects: first, the cloud computing mass data processing technology is introduced to provide strong support for large-scale WSN data processing; second, for the large-scale WSN itself, the innovative architecture suitable for the massive data processing is put forward, which can better integrate with cloud computing.

3 Method

3.1 Sensor network structure Map

There are N nodes in the WSN, each node has a unique identifier number, which has nothing to do with the location of the node. Sink is deployed at the center of the region. The level number is allocated from 1, the most inner node (nearest to Sink) has the smallest level number, and the outermost node has the largest level number.

Define GCP for the global control frame sent by Sink node, which can be received by all nodes in the Sink managed area. In network initialization, Sink broadcasts a GCP message with the area radius R, which contains the total number L of the ring network and the local communication reference radius $r_{com}$. The Sink node receives the local control frames sent by the neighbor nodes within the distance $r_{com}$, and the sensor nodes within the Sink node $r_{com}$ will become the first layer of nodes, that is, the virtual group head. Nodes will calculate the ring area limit value $b_k$ of the kth layer according to the GCP message and Formula (1).

$$b_k = \begin{cases} 
0, & k \leq 0 \\
\bar{b}_{k-1} + r_{com}, & \text{otherwise} 
\end{cases} \quad (1)$$

The parameter $\bar{b}_{k-1}$ represents the average distance from Sink node to the nodes in the k-1th layer, and its formula is shown in (2):

$$\bar{b}_{k-1} = \sqrt{\frac{b_{k-1}^2 - b_{k-2}^2}{2}}. \quad (2)$$

In the data acquisition stage, the node of the k+1 layer will transfer the sensor data to the nodes in the kth layer, the node density $\rho$ is usually a large value, and the sensor nodes are approximately uniform. Each node in the kth layer will receive approximately $\bar{T}_k$ data items from the k+1th layer node. $\bar{T}_k$ is described in Formula (3).
When all nodes have identified their own layer, they will enter the stage of identifying their neighbor nodes. When the neighbor node receives the request message and allows the connection, the acknowledgement request message will be sent to the source node, so that the communication links between the two sides will be established.

3.2 Data storage Map

The data storage Map used to segment the network vertically is mainly completed by Sink nodes. In other words, the Sink node first calculates the average density of each layer of nodes, and then divides the lattice according to the coverage of the nodes in each layer, so that the node density in each lattice is close. This means that in the same layer, the number of nodes in each lattice is approximately equal. The nodes in the lattice should be nodes in the neighbor tables of each node. If not, Sink should be notified to redistribute nodes between the lattice and neighbor lattice. If the SH node in $C_{ij}$ is selected, then Sink first calculates the distance between each node in $C_{ij}$ and then collects the residual energy information of the node, and finally calculates the possibility of each node becoming the memory head node according to the Formula (4).

$$P_{SH} = \frac{E_{res,k}}{d_k - b_t}.$$ 

(4)

4 Experiment simulation and discussion

4.1 Experimental simulation and discussion of SVC4WSN architecture

First of all, the multi-Sink mechanism in the sensor network layer is discussed by experiments. Assuming that there are totally 1000 sensor nodes randomly distributed in the monitoring area, the length of square is 100m and R is 50m, and there are a total of 4 nodes deployed in this area, then the implementation results based on Inter-Area Load Balancing algorithm is shown in Figure 1.

Figure 2 is the comparison of load balancing when the number of Sink nodes is 4. The horizontal coordinate shows the sensor node deployment number from 1000 to
10000, and the vertical coordinate suggest the deviation between 4 Sink nodes and average number of nodes before and after load balancing.

![Fig. 1. The implementation results based on Inter-Area Load Balancing algorithm](image1)

![Fig. 2. The comparison of deviation results of Sink nodes before and after load balancing](image2)
From the figure, it can be seen that, after load balancing, the deviation percentage between Sink node and average number of nodes is smaller than that before load balancing. And with the increase of the number of sensor nodes, that is, the density of nodes in the area becomes larger, the comparison of two deviation percentage lines shows good distance. It proves that Inter-Area Load Balancing algorithm has good extensibility in terms of load balancing effect between large-scale sensor networks.

Three ideal Sink nodes are deployed in the same area. It is supposed that the rectangle border length is 50m, the center of the circular area is three Sink nodes arranged in triangle and the distance is 5m, and they constitute an equilateral triangle topology. Ordinary sensor nodes are circular, and ideally, they are evenly partitioned into three sector regions by three Sink. Figure 3 shows the comparison of the number of nodes in the three Sink before and after using the Inter-Area Load Balancing algorithm. The total number of sensor nodes in the area is 900.

![Fig. 3. The comparison of the number of nodes in the three Sink before and after using the Inter-Area Load Balancing algorithm](http://www.i-joe.org)
Table 1. The parameters of quality of service for 5 services (obtained over the same period)

<table>
<thead>
<tr>
<th>Service provider</th>
<th>Response time</th>
<th>Throughput</th>
<th>Cost ratio</th>
</tr>
</thead>
<tbody>
<tr>
<td>S1</td>
<td>0.32</td>
<td>0.23</td>
<td>0.212</td>
</tr>
<tr>
<td>S2</td>
<td>0.24</td>
<td>0.30</td>
<td>0.288</td>
</tr>
<tr>
<td>S3</td>
<td>0.12</td>
<td>0.15</td>
<td>0.165</td>
</tr>
<tr>
<td>S4</td>
<td>0.21</td>
<td>0.08</td>
<td>0.233</td>
</tr>
<tr>
<td>S5</td>
<td>0.11</td>
<td>0.24</td>
<td>0.102</td>
</tr>
</tbody>
</table>

It can be found from the table that S1 takes the longest response time in the same time period, and S2 can get the maximum throughput on the virtual machine, but S2 spends the most virtual machine resources and S5 spend the least resource. Using the formula of computing service Rank value provided in the algorithm, the service Rank value and service sequencing shown in Table 2 can be obtained.

Table 2. The order of Rank value of 5 services (a=0.4, b=0.2, c=0.4)

<table>
<thead>
<tr>
<th>Service provider</th>
<th>Service Rank</th>
<th>Service order</th>
</tr>
</thead>
<tbody>
<tr>
<td>S1</td>
<td>0.2588</td>
<td>4</td>
</tr>
<tr>
<td>S2</td>
<td>0.2712</td>
<td>5</td>
</tr>
<tr>
<td>S3</td>
<td>0.144</td>
<td>2</td>
</tr>
<tr>
<td>S4</td>
<td>0.1932</td>
<td>3</td>
</tr>
<tr>
<td>S5</td>
<td>0.1328</td>
<td>1</td>
</tr>
</tbody>
</table>

In the case of a=c=0.4 and b=0.2, the service with the lowest Rank value is S5. At this time, the service Rank is more concerned with response time and resource cost, so a and c are valued larger than b. Selecting a smaller throughput weight b can reduce the effect of the throughput rate on the service Rank value. It is also what the virtual platform expects for the quality of service, that is, to achieve higher service throughput when the response time and the platform resource cost ratio is small, which is the basis of the Rank values of service S5 ranked No.1. If the response time threshold and throughput threshold of the virtual machine platform are set at 20%, then the response time and throughput rates of service S1 and S2 are higher than the set threshold. If the cost ratio of the two service platforms does not exceed the cost state value of the current virtual machine (VM) platform (this is a variable, depending on the number of services running on the VM platform simultaneously), then S1 and S2 will be transferred to a new VM platform; otherwise, the operation of services S1 and S2 will be suspended.

4.2 Experimental simulation and discussion of MDF4LWSN frame

In the Map of the MDF4LWSN framework network structure, the most important design goal is to obtain the optimal number of network layers and the unit communication radius of the corresponding nodes in the case of the minimum total energy consumption.
Figure 4 shows the relationship between the network optimal layer number $L$ and the node communication radius $r_{\text{com}}$ in the circular range of the radius $R=100\text{m}$ in the case of different node density $P$ (number of nodes contained in per square meter). The node density $P$ varies from 0.1 to 0.3, representing the dashed line at the top when $p=0.3$, and the dashed line at the bottom when $p=0.1$. It can be seen from the figure that with the decrease of node density $p$, the communication radius between nodes will gradually increase, and the corresponding optimal number of layers will decrease. This is because, in order to minimize the total energy consumed by the whole network, the possibility of communication between nodes is increased and the distance between the nodes is closer in the case of the number of node increasing. Therefore, the communication radius $r_{\text{com}}$ of the node will become smaller and the network optimal layer number $L$ will increase, and $L$ will be obtained when the optimal value $E_{\text{SUM}}$ is the minimum. The dotted lines with high $P$ values are above the dotted lines with low $p$ values, which shows that when the node communication radius is constant, the nodes need more layers to make the whole network less energy-consuming.

Consider the relationship between energy Reduce path and data storage Map. The design purpose of energy Reduce path design is to minimize the energy consumption between storage head (SH) nodes. In the area where radius $R=100\text{m}$, there are 3000 nodes in total, and the node density value is about 0.1. When the node communication radius $r_{\text{com}}$ equals 2m, the optimal layer number of the network is 40 layers. The number of bits per second of each node is 100bit/s. In the case of time $T=10\text{s}$, Figure 5 shows the effect of the data storage Map on the energy consumed by the energy Reduce path.
As can be seen from the figure that, the energy consumed by the whole network under the action of the energy Reduce path has a significant downward trend when there is a data storage Map. This is because in the data storage Map, a space time map is taken to the node, so that the sub nodes belonging to the same SH node only collect the data belonging to their own time slot. Although the burden of the SH node is increased, the energy of the whole network is saved by the existence of the backbone path. From the two curves, it is seen that, as the number of nodes increases, the energy consumption is greater. But in general, the increase of energy consumption is proportional to the increase of the number of nodes. This shows that the number of nodes on the backbone path is not significantly changed in the construction of the energy Reduce path, and the most energy consumption is the node collecting data and the SG node.

Finally, the performance of the Reduce path in the MDF4LWSN framework is discussed, and it is compared with the three protocols of LEACH, PEGASIS and PEDAP. Figure 6 shows the performance of data acquisition delay under four different protocols.

The time slot number of data acquisition based on the time delay Reduce path is the smallest. The reason is that in this case, the number of time slots is the logarithm of the number of nodes, and with the increase of the number of nodes, the number of time slot is maintained in a stable state. PEDAP is followed by it in terms of the delay performance, which is larger than MDF4LWSN, and they are tree topology. But PEDAP cannot guarantee that the parent node does not have to wait for the time slot completion of the brother node when the data is received. Therefore, the tree topology constructed with the shortest path algorithm does not have the optimal transmission delay in the PEDAP. The number of time slots collected by PEGASIS is much higher than that of the previous two algorithms, because in the link, it takes more time to transfer data from the terminal node to the middle head node in the link. The delay performance of LEACH is the worst, and on average, it is almost proportional to the number of nodes.
5 Conclusion

On the basis of the comprehensive analysis of WSN and cloud computing technology, some suggestions on the architecture of WSN based on cloud computing and the large-scale data processing scheme are put forward. The large-scale WSN architecture based on cloud computing and the mass data processing framework in the WSN is studied.

The effect of multi-Sink mechanism on the load balancing of WSN nodes is discussed through simulation experiments, and the comparison of the results before and after load balancing is carried out. In the Map part of the network structure, the relationship between the optimal number of nodes and the communication radius of nodes under different node densities is verified by simulation. The experiment simulates the influence of data storage Map on network energy consumption. Finally, the energy and delay Reduce paths are discussed and compared with the three protocols of LEACH, PEGASIS and PEDAP in simulation.

6 References


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